

## RelationalAD

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### **Rel**AD: What is the Rel language?

- Declarative				
- Multipurpose				
- Logic				
- Database queries				
- Linear algebra				1
- Tensor computation				
- Machine learning				
- Feature extraction				
- Modeling, inference, prediction				
- Mathematical Optimization				
- Statistics				
<ul> <li>Probabilistic programming</li> </ul>				
- See <u>relational.ai</u>				
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### RelAD: Rel Core Syntax

- A (vast) generalization of <i>Datalog</i> with agg/neg	
- A Rel program is a collection of rules	
$-$ def Q(x, y,) = $\phi(x, y,)$	
- A Formula φ(x , y,) defines a <i>relation</i> over vars {x, y,}	
- Each formula $\varphi(x, y,)$ could be	
- A materialized atom, e.g. R(x, y,)	
- A native, e.g. $x + y = z$ or $x > y$	
- A conjunction/disjunction of formulas, e.g. $\psi_1(x,) \wedge \psi_2(x,)$	(2000)
- A negation, e.g. $\neg \psi(x, y,)$	
- ∃ or ∀, e.g. ∃ z : ψ(x, y, z,)	
- A sum/reduce, e.g. sum[z, t: ψ (x, z, t,)](y)	
- An FFI, e.g. some-external-function( $\psi$ )	
See <u>docs.relational.ai/rel/primer/basic-syntax</u>	
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### Rel<mark>AD: Rel Examples</mark>

Triangle counting in a graph E def Q = count[a, b, c: E(a, b) and E(b, c) and E(c, a)] Matrix multiplication C = ABdef C(i, j, v) = sum[ k, v1, v2, v3: A(i, k, v1) and B(k, j, v2) and v1\*v2=v3](v) def C[i, j] = sum[k: A[i, k] \* B[k, j]]  $J = \|Ax - b\|_2^2$ def J = sum[i : (sum[j : A[i, j] \* x[j]] - b[i])^2] 5 

### **RelAD**: Differentiation (What We Learned in College)

$$egin{aligned} &rac{\partial eta^ op X^ op X eta c}{\partial X} = X(eta c^ op + eta b^ op) \ &rac{\partial \log |A|}{\partial A} = (A^{-1})^ op \ &rac{\partial \log |A|}{\partial A} = (A^{-1})^ op \ &rac{\partial \operatorname{tr}(BA)}{\partial A} = B^ op \ &rac{\partial eta^ op X^ op D X eta c}{\partial X} = D^ op X eta c^ op + D X eta b^ op \ &rac{\partial (Xb+c)^ op D (Xb+c)}{\partial X} = (D+D^ op) (Xb+c) b^ op \end{aligned}$$



### **RelAD**: What is Differentiation Used For?

- Most of modern machine learning

- Traditional optimization







### **Rel<mark>AD</mark>: Automatic Differentiation**

- Input: a **program** computing a function *f*
- Output: a **program** computing *df*

#### What is a program?

- A neural network
- Imperative program in C++, Haskell, etc
- A **Rel program**! (even with recursion)

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### **Some Existing AutoDiff Frameworks**

Tapcar Flow (autodiff)	
- TensorFlow ( <u>autodiff</u> )	
- PyTorch ( <u>autograd</u> )	
- Tyrorch ( <u>autograu</u> )	
- NumPy ( <u>IAX</u> )	
- <u>Geno</u>	
They operate on <b>Einsum notations</b> to construct complex functions of tensors	
They operate of Ensuin notations to construct complex functions of tensors	



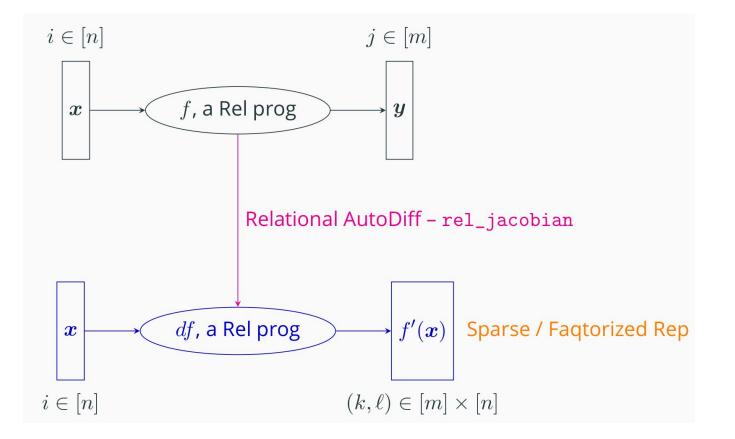
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### **Digression: Einstein Notation**

Existing AutoDiff frameworks operate on (network of) *Einsum rules*, e.g.  $U_{i,j,k} = \sum_{l.m} R_{i,l,m} \cdot S_{j,k,l} \cdot T_{i,k,l}$ def U[i, j, k] = sum[1 m v : v = R[i,1,m] \* S[j,k,1] \* T[i,k,1]] (For a gentle intro, see "<u>Einsum is all you need</u>") Rel is **much** more general, e.g. def U[i, j, k] = sum[1 m v : v = R[i,1,m] \* S[j,k,1] \* T[i,k,1] and  $exists(x : i^2 + x \le k and x > 10)$ • • • • • Our keys are of arbitrary types - Tensors can be (very) sparse - Additional logic is arbitrary - Worst-case optimal join + semantic optimization + IVM 

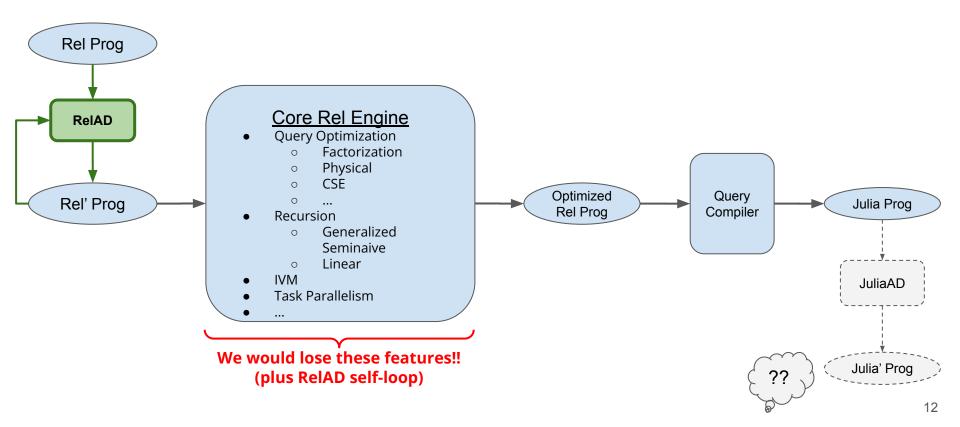


### RelAD: Relational AutoDiff





### Why RelAD, not JuliaAD?





### **RelAD: Example**

```
Consider the Rel program
                                                                     J = x^T A x^{-1}
        def J = sum[i j : x[i] * A[i, j] * x[j]]
   We need to give an extra hint to specify how to interpret it as a function

abla = rac{\partial J}{\partial x}
         def \nabla = jacobian[J, x]
   This says that
     - the above program defines a function f: x \rightarrow J, and
                                                                                               • • we are interested in f', which we now call \nabla.
```



### RelAD: Example (Cont.)

```
RelAD rewrites this program into
                                                                           J = x^T A \overline{x} .
def J = sum[i j : x[i] * A[i, j] * x[j]]

abla_1 = Ax
def \nabla 1[i] = sum[j v : x[i] = and v = A[i, j] * x[j]]

abla_2 = A^T x^{-1}
def \nabla 2[j] = sum[i v : v = x[i] * A[i, j] and x[j] = ]
                                                                 \nabla = \nabla_1 + \nabla_2
def \nabla[i] = merge sum[\nabla1[i], \nabla2[i]]
                                                                                     14
```



### RelAD: Example (Cont.)

Now take this new program and add the hint

```
def H = jacobian[\nabla, x]
```

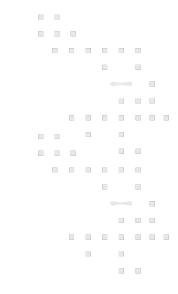
The above says

- interpret the new program as another function  $g: x \rightarrow \nabla$ , and

 $H = \frac{\partial \nabla}{\partial x}$ 

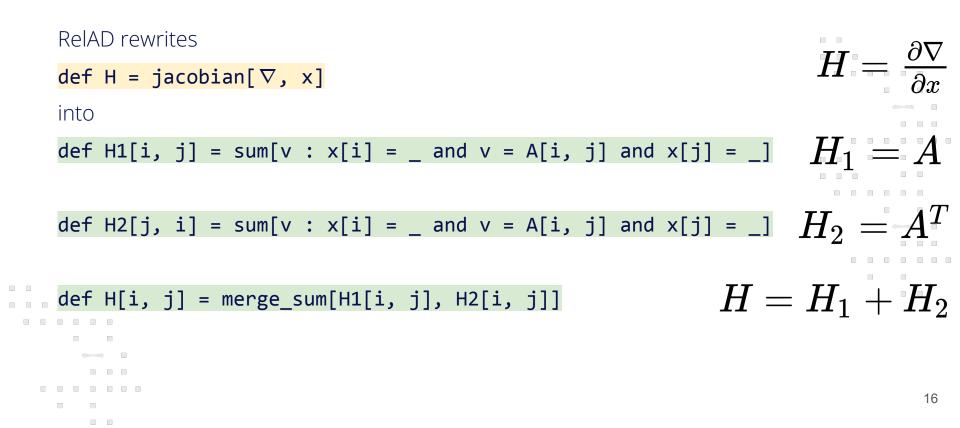
- compute g' which we now call H







### RelAD: Example (Cont.)



### Interface: "jacobian" Higher-order Native

- Given a Rel program P defining (among other things) two relations A and B where
  - A[k<sub>1</sub>, k<sub>2</sub>, ..., k<sub>m</sub>] = v has m >= 0 keys and one value v whose type is Float
  - **B**[**I**<sub>1</sub>, **I**<sub>2</sub>, ..., **I**<sub>n</sub>] = **w** has **n** >= **0** keys and one value **w** whose type is **Float**
  - **B** may depend on **A** in any way: directly or indirectly through chains of other relations in **P**
- We can use the higher-order native **jacobian** to define a new relation **C** 
  - def C = jacobian[B, A]

- C[l<sub>1</sub>, l<sub>2</sub>, ..., l<sub>n</sub>, k<sub>1</sub>, k<sub>2</sub>, ..., k<sub>m</sub>] = t has n+m keys and one value t whose type is Float
- $C[l_1, l_2, ..., l_n, k_1, k_2, ..., k_m] := \partial B[l_1, l_2, ..., l_n] / \partial A[k_1, k_2, ..., k_m]$

```
- RelAD later desugars jacobian into lower-order Rel
```

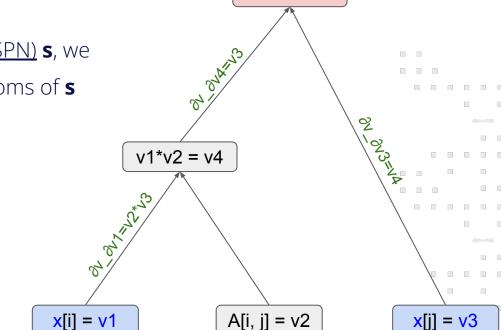


#### relational<u>A</u>

### How does it work?

To derive <u>a single SumProductNode (SPN)</u> **s**, we

- analyze dependencies among atoms of s
- construct a dependency DAG
- do backpropagation

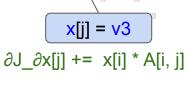


 $\partial J \partial x[i] += A[i,j] * x[j]$ 

v4\*v3 = v

Example:

def J = sum[i j : x[i] \* A[i, j] \* x[j]]  $J = x^T A x$ def  $\nabla$  = jacobian[J, x]  $\nabla = \frac{\partial J}{\partial x}$ 





### How does it work?

To derive <u>an entire Rel program</u> <b>P</b> consisting of many SPNs, we	
<ul> <li>analyze dependencies among SPNs of P</li> </ul>	
<ul> <li>construct a dependency DAG</li> </ul>	(an-ra)
- do <b>backpropagation</b>	
- More generally, Jacobian accumulation	
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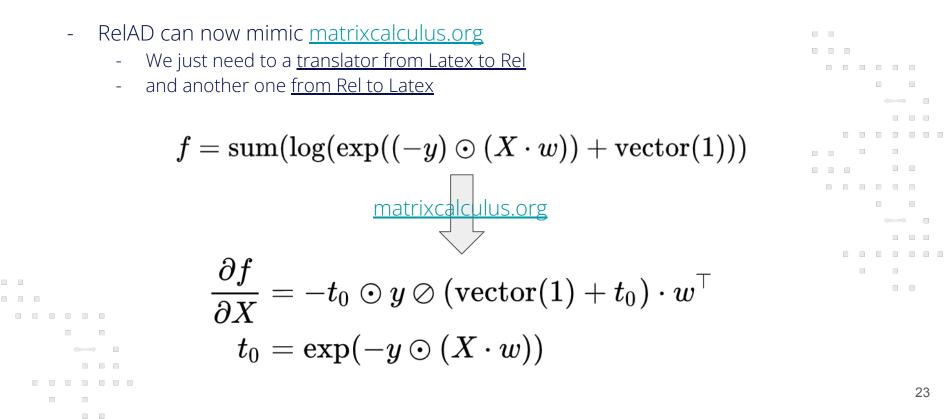
(2000) 

### **More Matrix Calculus Examples**

#### From our test suite:

$$\nabla (Bx+b)^{\top} C(Dx+d) = B^{\top} C(Dx+d) + D^{\top} C^{\top} (Bx+b)$$
$$\frac{\partial b^{\top} X^{\top} X c}{\partial X} = X(bc^{\top} + cb^{\top})$$
$$\frac{\partial \operatorname{tr}(BA)}{\partial A} = B^{\top}$$
$$\frac{\partial b^{\top} X^{\top} DX c}{\partial X} = D^{\top} X bc^{\top} + DX cb^{\top}$$
$$\frac{\partial (Xb+c)^{\top} D(Xb+c)}{\partial X} = (D+D^{\top})(Xb+c)b^{\top}$$

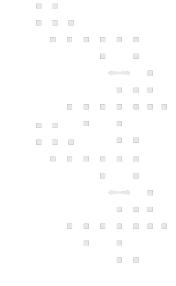
### **More Matrix Calculus Examples**





### **Latex** $\Rightarrow$ **Rel** (currently missing)

 $f = \mathrm{sum}(\log(\exp((-y)\odot(X\cdot w)) + \mathrm{vector}(1)))$ 



```
def z[i] = sum[j : -y[i] * X[i, j] * w[j]]
  def f = sum[i : natural_log[natural_exp[z[i]] + 1.0]]
  def ∇ = jacobian[f, X]
```



### $\mathbf{Rel} \Rightarrow \mathbf{Rel} \qquad (\mathbf{RelAD})$

def z[i] = sum[j : -y[i] \* X[i, j] \* w[j]]def f = sum[i : natural\_log[natural\_exp[z[i]] + 1.0]] def  $\nabla$  = jacobian[f, X]

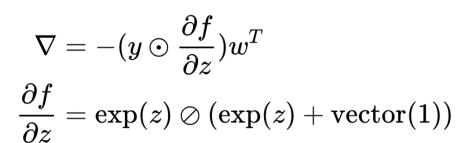
def  $\nabla$ [i, j] = v : X[i, j] = \_ and v = -y[i] \* w[j] \*  $\partial$ f\_ $\partial$ z[i] def  $\partial$ f\_ $\partial$ z[i] = natural\_exp[z[i]] / (natural\_exp[z[i]] + 1)

RelAD



### **Rel** $\Rightarrow$ **Latex** (currently missing)

def  $\nabla$ [i, j] = v : X[i, j] = \_ and v = -y[i] \* w[j] \*  $\partial$ f\_ $\partial$ z[i] def  $\partial$ f\_ $\partial$ z[i] = natural\_exp[z[i]] / (natural\_exp[z[i]] + 1)



### **ML Example: Multi-layer Neural Network**

Consider a neural network with two layers, ReLU activations, and multiple outputs:

-	old x is the input vector,	<b>t</b> is the target vector		_	
-	$W_1$ is the weight matrix for th	ne 1st layer			
	W, is the weight matrix for th				
	Recall $\operatorname{ReLU}(x) := \max(x)$				
		, •)			
Th	e network works as follows:	$y_1=W_1x$			
		$z_1 = \operatorname{ReLU}(y_1)$			
		$y_2=W_2z_1$			
		$z_2 = \operatorname{ReLU}(y_2)$			
		$J = \ t - z_2\ _2^2$			
				2	27
				_	

### **ML Example: Multi-layer Neural Network (cont.)**

Input Rel Program: $y_1 = W_1 x$ def y1[i] = sum[j : W1[i, j] \* x[j]]

 $z_1 = \operatorname{ReLU}(y_1)$ def z1(i, v) = y1[i] = v and v >= 0 def z1(i, v)=exists(u : y1[i]=u and u<0 and v=0)

 $y_2 = W_2 z_1$ def y2[i] = sum[j : W2[i, j] \* z1[j]]

 $z_2 = \operatorname{ReLU}(y_2)$ def z2(i, v) = y2[i] = v and v >= 0 def z2(i, v)=exists(u : y2[i]=u and u<0 and v=0)</pre>  $J = \|t - z_2\|_2^2$ def J = sum[i : (t[i] - z2[i]) ^ 2]  $\nabla_{W1} = \frac{\partial J}{\partial W_1}$ def ▽ W1 = jacobian[J, W1]  $\nabla_{W2} = \frac{\partial J}{\partial W_2}$ def ∇ W2 = jacobian[J, W2]



### **ML Example: Multi-layer Neural Network (cont.)**

RelAD output:

 $rac{\partial J}{\partial z_2}=2(z_2-t)$ def  $\partial J \partial z^{[i]} = 2 * (z^{[i]} - t^{[i]})$  $rac{\partial J}{\partial y_2} = rac{\partial J}{\partial z_2} \odot 1_{y_2 \geq 0}$  def  $\partial J \partial y_2[i] = sum[v: y_2[i] >= 0 and <math>\partial J \partial z_2[i] = v]$  $abla_{W2} = rac{\partial J}{\partial u_2} z_1^T$ def  $\nabla_W2[i, j] = sum[v: W2[i, j] = _ and v = \partial J_\partial y2[i] * z1[j]]$  $rac{\partial J}{\partial z_1} = W_2^T rac{\partial J}{\partial y_2}$ def  $\partial J_{\partial z1[j]} = sum[i v: z1[j] = _ and v = W2[i, j] * <math>\partial J_{\partial y2[i]}$  $rac{\partial J}{\partial y_1} = rac{\partial J}{\partial z_1} \odot 1_{y_1 \geq 0}$  . def  $\partial J_{\partial y1}[i] = sum[v: y1[i] >= 0 and <math>\partial J_{\partial z1}[i] = v$ ]  $abla_{W1} = rac{\partial J}{\partial u_1} x^T$ def  $\nabla_W1[i, j] = sum[v: W1[i, j] = _ and v = \partial_J_\partial y1[i] * x[j]]$ 29



### **Recursive Example**

$$rac{\partial \log(|\det(A)|)}{\partial A} = (A^{-1})^T$$

- The determinant is not directly available in Rel
- But it can be computed **recursively** using a Gram-Schmidt process







### **Recursive Example:** (cont.)

- Given an (n X n)-matrix A = [A<sub>1</sub>, A<sub>2</sub>, ..., A<sub>n</sub>]  
- Compute orthogonal vectors O = [O<sub>1</sub>, O<sub>2</sub>, ..., O<sub>n</sub>]  

$$O_j = A_j - \sum_{k < j} \frac{A_j^T O_k}{||O_k||_2^2} O_k$$
def O(i, j, v) = (j = 1 and v = A[i, 1])  
def O(i, j, v) = (v = A[i, j] - sum[k s :  
k < j and s = prod\_A\_0[j, k] / sqr\_norm\_0[k] \* 0[i, k]])  
def prod\_A\_0[j, k] = sum[1 : A[1, j] \* 0[1, k]]  
def sqr\_norm\_0[k] = sum[1 : 0[1, k] ^ 2]



### **Recursive Example:** (cont.)

- By determinant properties, det(A) = det(O)
- Because O is an orthogonal basis

$$|\det(O)| = ||O_1||_2 imes ||O_2||_2 imes \cdots imes ||O_n||_2$$
  
 $\log |\det(O)| = \log ||O_1||_2 + \log ||O_2||_2 + \cdots + \log ||O_n||_2$ 

def log\_det = sum[k : natural\_log[sqr\_norm\_0[k]] / 2]
def ∇ = jacobian[log\_det, A]





### **Example: Linear regression with Gradient-descent**

<ul> <li>Input         <ul> <li>(N×d)-matrix X where each row is a data point</li> </ul> </li> </ul>			
- N-vector t of corresponding target responses			
		(j	
- Output			
- d-vector w <sup>*</sup> of optimal model parameters			
$w^* = \mathrm{argmin}_w \ Xw - t\ _2^2$			
w = g = w = w = v = v = 2			
	_		
		_	
			34

### **Example: Linear regression with Gradient-descent**

```
def MAX K = 10000
                        // maximum number of iterations
     def \alpha = 0.01 // learning rate (fixed)
                                                                          w_0 \leftarrow 0
     def w(k, i, v) = k = 0 and range(1, d, 1, i) and v = 0.0
                                                                        y_k \leftarrow X w_k
     def y[k, i] = sum[j : X[i, j] * w[k, j]]
                                                                                           J_k \leftarrow \|y_k - t\|_2^2
     def J[k] = sum[i : (y[k, i] - t[i]) ^ 2]

abla_k \leftarrow rac{\partial J_k}{\partial w_k}
     def \nabla = jacobian[J, w]
                                                                                                    def w(k, i, v) = k <= MAX K and
                                                                    w_k \leftarrow w_{k-1} - lpha 
abla_{k-1}
   v = w[k-1, i] - \alpha * \nabla[k-1, k-1, i]
                                                                                                       35
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```



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### **Incorporating ICs: Key alignment**

Consider the program		
def J = sum[i : b[i] * x[i]] $J = b^T x$		
def $\nabla$ = jacobian[J, x] $\nabla - \frac{\partial J}{\partial J} - h$		
def $ abla$ = jacobian[], x] $ abla = b$		
Currently ReIAD produces the following rule for $ abla$ :		
def $\nabla$ [i] = v : v = b[i] and x[i] = _		
The <b>extra part</b> is due to the fact that ReIAD doesn't know that the keys of <b>b</b> an aligned	d x are	
<ul> <li>Even if there was an IC saying they are, our RelAD currently has no mecha</li> <li>utilize it</li> </ul>	nism to	
		37

### **Incorporating ICs: Matrix Symmetry**

Consider the example from before  

$$\begin{aligned}
\text{def J = sum[i j : x[i] * A[i, j] * x[j]} & J = x^T A x \\
\text{If A was symmetric, then} & \nabla = 2Ax
\end{aligned}$$
However currently our RelAD output has no mechanism to utilize this symmetry (even if it was encoded as an IC)  

$$\begin{aligned}
\text{def } \nabla 1[i] = \text{sum[j } v : x[i] = \_ \text{and } v = A[i, j] * x[j]] & \nabla_1 = Ax \\
\text{def } \nabla 2[j] = \text{sum[i } v : v = x[i] * A[i, j] \text{ and } x[j] = \_] & \nabla_2 = A^T x \\
\text{def } \nabla [i] = \text{merge}_{\text{sum}} [\nabla 1[i], \nabla 2[i]] & \nabla = \nabla_1 + \nabla_2
\end{aligned}$$



### **Multi-headed Rules**

Consider the example	
$// J = x^{T}Ax$	
def J = sum[i j : x[i] * A[i, j] * x[j]]	
def $\nabla$ = jacobian[J, x]	
Instead of having two rules for $ abla$	
$\nabla$ [i] += A[i, j] * x[j]	
$\nabla$ [j] += A[i, j] * x[i]	
it is faster to have one with two heads	
$\nabla$ [i] += v2*v3, $\nabla$ [j] + = v1*v2 $\leftarrow$ x[i] = v1, A[i, j] = v2, x[j]	= v3
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### **Other Issues**

-	Usage				
	- Integration with ML work				
	<ul> <li>Integration into existential second order (ESO) work (for optimization)</li> </ul>				
-	Performance				
	- Dense-tensor support				
	- Mapping to BLAS / GPUs /				
	- XY-Stratification to speed up GD (Joint work with Amir Shaikhha)				
				0	







### **Some References**

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# Thank You!

Any Questions/Comments?

