



Capturing Factorized Provenance (FDB 2022)

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Outline

Overview & Motivation

- Provenance Graph Model
- Capturing Provenance
- Factorization
- Experiments
- Conclusions & Future work



Debugging Datalog

Datalog has seen a resurgence in academia & industry

- Programming distributed systems
- Complex data-intensive computations
- Analytics & ML over query results
- Urgent need for debugging Datalog programs
 - Why did my program produce this unexpected result?
 - Why did my program not produce this expected result?
 - Which rules are responsible for deriving this result?
 - Why did this rule derivation fail?
 - Would deleting & inserting a tuple change the result?





PUG + PUGS



Efficient capture and summarization for why and whynot provenance through a provenance model developed

for queries with negation





PUG + PUGS



Efficient capture and **summarization** for **why and whynot** provenance through a **provenance model developed**

for queries with negation

This is FDB!



PUG utilizes a flat-relational encoding of provenance graphs which essentially corresponds to a factorized representation of provenance



ILLINOIS INSTITUTE Schism between Why / Why-not

 How the rules of a program did derive / failed to derive an existing / missing output from the input data





Unifying Why / Why-not

• Why (provenance) and why-not (missing answers) have been mostly treated in isolation

- Why and why-not questions can be reduced to each other for a query language **L** if ...
 - for any query Q, its complement Q^c is in L

$$t \notin Q(D) \equiv t \in Q^C(D)$$

$$Q(A) := R(A,B)$$

 $Q^{C}(A) := R(A,B)$, adom(A), adom(B)



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Requirements

Syntax-driven

• Provenance structure aligns with program structure

Compatible with well-established provenance models

- Provenance polynomials for positive queries [1]
- Dual polynomials [2]

Build-in support for sharing common subexpressions

[1] T. Green, G. Karvounarakis, and V. Tannen. Provenance semirings. In PODS, pages 31–40, 2007.

[2] E.Grädel and V.Tannen.Semiring provenance for first-order model checking. arXiv preprint arXiv:1712.01980, 2017.



tuple

rule

goal

tuple





























• Equivalent model (annotated rule derivations)



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Extracting Provenance Polynomials





 $r_{2}(g_{1}(p) \cdot g_{2}(p) \cdot g_{3}(p)) + r_{2}(g_{1}(q) \cdot g_{2}(r) \cdot g_{3}(p)) + r_{2}(g_{1}(p) \cdot g_{2}(q) \cdot g_{3}(r))$





Factorized Provenance ILLINOIS INSTITUTE Graphs







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Computing Provenance

Problem definition

input: (missing) answers of interest
output: relevant subgraph for all answers of interest

Instrumentation of the input Datalog program

- Use firing rules to capture rule derivations
- Goal-oriented approach
 - Limit the provenance based on the user question
 - Efficient bottom-up evaluation using relational engines
 - Compiling rewritten program into SQL



Overview

Rewriting steps

- **Unifying** the program with specification of outputs of interest
- Statically annotate program to indicate interest in success / failure
- Creating firing rules
- Checking connectivity with joins
- **Computing** the **edge** relation of the provenance graph

Example query and provenance question

$$r_1: Q(X, Y): -\operatorname{Train}(X, Z), \operatorname{Train}(Z, Y), \neg \operatorname{Train}(X, Y)$$

WHY Q(n,s)?



Computing Provenance

Rewriting steps

- **Unifying** the program with specification of outputs of interest
- Statically annotate program to indicate interest in success / failure
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- **Computing** the **edge** relation of the provenance graph



Step 1: Unifying the program with the question

- Limit computation by binding variables to constants
- Propagating constants in the question throughout the program

Example

Binding variables: X = n (New York) and Y = s (Seattle)

 $r_1: Q(X, Y): -\operatorname{Train}(X, Z), \operatorname{Train}(Z, Y), \neg \operatorname{Train}(X, Y)$





Step 2: Annotating rule with success / failure

- To determine the state of nodes
- The question provides information about success / failure
 True = Success / Exist
 False = Failed / Not Exist
- Propagate annotations throughout the unified program

Example

▷ WHY Q(n,s) implies only successful goals need to be captured



Step 3: Create firing rules

 Capture successful / failed derivations for the variable binding

Example

- Create new head by adding the variable Z and use firing versions
- Invert the annotation T for the goal ¬T(n,s) in the firing version

$$r_{1}^{(X=n,Y=s),T}: \boxed{\mathbb{Q}(n,s)^{T}}:- \underbrace{\mathbf{T}(n,Z)^{T}}_{\mathbf{T}(Z,s)^{T}}, \underbrace{\mathbf{T}(Z,s)^{T}}_{\mathbf{T}(n,s)^{T}} \xrightarrow{\mathbf{F}_{\mathbf{r}_{1},\mathbf{T}}(n,s,Z):- \underbrace{\mathbf{F}_{\mathbf{r}_{1},\mathbf{T}}(n,z)}_{\mathbf{F}_{\mathbf{T},\mathbf{T}}(n,Z),\mathbf{F}_{\mathbf{T},\mathbf{T}}(Z,s)}, \underbrace{\mathbf{F}_{\mathbf{T},\mathbf{F}}(n,s)}_{\mathbf{F}_{\mathbf{T},\mathbf{T}}(Z,s):- \underbrace{\mathbf{T}(n,Z)}_{\mathbf{F}_{\mathbf{T},\mathbf{T}}(Z,s),\mathbf{F}_{\mathbf{T},\mathbf{T}}(Z,s)}$$



Step 4: Connectivity joins

- Firing rules are not sufficient to determine which subgraphs of the provenance explain the outputs of interest
- Filter derivations by checking whether connectivity
- Check connectivity from the question node one hop at a time

Example

- No guarantees for the nodes in the red box
- Tuple node T(n,c) is only connected iff T(c,s) exists



$$\begin{split} \mathbf{F}_{\mathbf{Q},\mathbf{T}}(n,s) &:= \mathbf{F}_{\mathbf{r}_{1},\mathbf{T}}(n,s,Z) \\ \mathbf{F}_{\mathbf{r}_{1},\mathbf{T}}(n,s,Z) &:= \mathbf{F}_{\mathbf{T},\mathbf{T}}(n,Z), \mathbf{F}_{\mathbf{T},\mathbf{T}}(Z,s), \mathbf{F}_{\mathbf{T},\mathbf{F}}(n,s) \\ \hline \mathbf{F}\mathbf{C}_{\mathbf{r}_{2},\mathbf{r}_{1}^{1},\mathbf{T}}(n,Z) &:= \mathbf{T}(n,Z), \mathbf{F}_{\mathbf{r}_{1},\mathbf{T}}(n,s,Z) \\ \mathbf{F}\mathbf{C}_{\mathbf{r}_{2},\mathbf{r}_{1}^{2},\mathbf{T}}(Z,s) &:= \mathbf{T}(Z,s), \mathbf{F}_{\mathbf{r}_{1},\mathbf{T}}(n,s,Z) \\ \mathbf{F}_{\mathbf{T},\mathbf{F}}(n,s) &:= \neg \mathbf{T}(n,s) \end{split}$$



Step 5: Computing provenance subgraph edge relation

- Create edges for the provenance graph (explanation)
- Generate rules for the edge relation based on the rule binding information
- Use node identifier $\int_{(1)}^{(2)} (n, s, Z)$
- Type of the node, assignments to constants, success/failure state
- Each rule corresponds to a pattern in the graph



- Provenance graph structure -

$\texttt{edge}(f_Q^T(n,s),f_{r_1}^T(n,s,Z))\!:=\!\texttt{F}_{\texttt{r}_1,\texttt{T}}(n,s,Z)$	Z)
$edge(f_{r_1}^T(n, s, Z), f_{g_1^1}^T(n, Z)) := F_{r_1, T}(n, s, T)$	Z)
$\texttt{edge}(f_{g_1^1}^T(n,Z),f_T^T(n,Z))\!:=\!\texttt{F}_{\texttt{r}_1,\texttt{T}}(n,s,I)$	Z)
$edge(f_{g_1^3}^T(n,s), f_T^F(n,s)) := F_{\mathtt{r}_1, \mathtt{T}}(n,s, t)$	Z)

Example (partial) rules deriving the edge relation



Implementation

PUG (Provenance Unification through Graphs) architecture

- Extension of GProM supporting Datalog provenance
- GProM is a SQL+X to SQL optimizing compiler
- Relational algebra as IR







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Our provenance graphs can encode factorized polynomials

- What factorization we get is determined by the structure of the program
- Q() := S(X), U(X,Y)
- Q() := T(X), U(X,Y)





Utilizing work on factorization



Approach

- Determine worst-case optimal d-tree factorization
- Rewrite input query to produce this factorization
- Apply capture rewriting

D. Olteanu and J. Závodny`. Size bounds for factorised repre- sentations of query results. ACM Transactions on Database Systems (TODS), 40(1):2, 2015.

 $r_3: \mathsf{Q}_{2\mathrm{hop}}(X):-\mathrm{H}(Y,L_1,Z), \mathrm{H}(Z,L_2,X)$



Utilizing work on factorization



Approach

- Determine worst-case optimal d-tree factorization
- Rewrite input query to produce this factorization
- Apply capture rewriting

$$\begin{array}{c} r_{5}: \mathtt{Q}_{2\mathrm{hop}}():- \mathtt{Q}_{\mathtt{L}_{1}}(Z), \mathtt{Q}_{\mathtt{L}_{2}}(Z) \\ r_{5'}: \mathtt{Q}_{\mathtt{L}_{1}}(Z):- \mathtt{H}(Y, L_{1}, Z) \\ r_{5''}: \mathtt{Q}_{\mathtt{L}_{2}}(Z):- \mathtt{H}(Z, L_{2}, d) \\ & \stackrel{\{\} \ Z}{\underset{\{Z\} \ L_{1} \ L_{2} \ \{Z\}}{} \\ \end{array}$$

 $\{Z, L_1\} Y$

Relation H

\mathbf{S}	${f L}$	\mathbf{E}	
a	l_1	с	s_1
a	l_2	с	s_2
b	l_3	с	t_1
b	l_4	с	t_2
с	l_5	d	u_1
С	l_6	d	u_2





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• **TPC-H** suppCust(N):-SUPPLIER(A,B,C,N,D,E,F), CUSTOMER(G,H,I,N,J,K,L,M)





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Conclusions

Provenance graph model

- Provenance structure aligns with program structure
- Compatible with well-established provenance models
- Provenance polynomials for positive queries + dual polynomials
- Build-in support for sharing common subexpressions
- Flat relational encoding as edge relation

Capturing Provenance

- Incorporate user's provenance interest into the capture query
- Filter successful / failed assignments upfront (static analysis)
- Output is a query returning the edge relation of the graph



Conclusions

Factorization

• Program structure determines factorization

Approximate Summarization of Why-not Provenance

• Use patterns to summarize provenance

• Use sampling to generate such summaries for very large why-not provenance graphs



Future work

Expressiveness

• Support aggregation and recursion

Efficiency

- Leverage factorized DB techniques for aggregates?
- Factorizing missing answers (complement representations)?

Going beyond SQL / Datalog as the target language

• What specialized algorithms & data structures would be beneficial?



Questions?







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PUG: https://github.com/IITDBGroup/PUG





Airbnb (bed and breakfast)

Listing (input)

\mathbf{Id}	Name	Ptype	Rtype	NGroup	Neighbor
8403	central place	apt	shared	queen anne	east
9211	plum	apt	entire	ballard	adams
2445	cozy homebase	house	private	queen anne	west
8575	near SpaceNeedle	apt	shared	queen anne	lower
4947	seattle couch	condo	shared	$\operatorname{downtown}$	first hill
2332	modern view	house	entire	queen anne	west

\mathbf{A} vaila	ability (input)	
\mathbf{Id}	Date	Price
9211	2016-11-09	130
2445	2016 - 11 - 09	45
2332	2016-11-09	350
4947	2016-11-10	40

r₁: AL(N,R) :- L(I, N, T, R,queen anne, E), A(I, 2016-11-09, P)

"What are avaliable listings and the room types in Queen Anne on Nov 9th, 2016?"







All derivations

Single derivation

Why-not [Huang et al. VLDB 2008] Artemis [Herschel et al. VLDB 2009] Y! [Wu et al. SIGCOMM 2014]

Not scalable

Not comprehensive



\mathbf{L} istin	g (input)				
\mathbf{Id}	Name	Ptype	Rtype	NGroup	Neighbor
8403	central place	apt	shared	queen anne	east
9211	plum	apt	entire	ballard	adams
2445	cozy homebase	house	private	queen anne	west
8575	near SpaceNeedle	apt	shared	queen anne	lower
4947	seattle couch	condo	shared	downtown	first hill
2332	modern view	house	entire	queen anne	west

\mathbf{A} vaila	bility (input)	
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9211	2016-11-09	130
2445	2016-11-09	45
2332	2016-11-09	350
4947	2016-11-10	40

Why no shared room exists?

r₁: AL(N,R) :- L[I,N,T]R queen anne, E), A[I, 2016-11-09,P]

Attribute	Id	Name	Ptype	Rtype	NGroup	Neighbor	Date	Price	- 2160
#Distinct Values	6	6	3	3	3	5	2	4	= 2100



~15·10²⁰ derivations over full dataset (~ 1.4M)





Listing (input)

\mathbf{Id}	Name	Ptype	Rtype	NGroup	Neighbor
8403	central place	apt	shared	queen anne	east

\mathbf{A} vaila	ability (input)	
\mathbf{Id}	Date	Price

r₁: AL(N,R) :- L(I, N, T, R,queen anne, E), A(I, 2016-11-09, P)



The listing 'central place' has a shared room which is not available at \$130 AL(central place, shared) r_1 (central place, shared, 8403, apt, east, 130) $g_1^2(8403, 2016-11-09, 130)$ A(8403, 2016-11-09, 130)



.. /.

Motivation (Example)

Listing (input)

\mathbf{Id}	Name	Ptype	Rtype	NGroup	Neighbor
8403	central place	apt	shared	queen anne	east
9211	plum	apt	entire	ballard	adams
2445	cozy homebase	house	private	queen anne	west
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\mathbf{Id}	Date	Price
9211	2016-11-09	130
2445	2016-11-09	45
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r₁: AL(N,R) :- L(I, N, T, R,queen anne, E), A(I, 2016-11-09, P)



" All shared rooms of apartments in Queen Anne are not available at any price on Nov 9th, 2016 "





Goals

- **Concise** (small size of explanations)
- Complete (covering all provenance)
- Informative (providing new insights)
- Challenge
 - Fullfilling all 3 elements at the same time
 - Computing summaries using full why-not provenance



- Computing top-k summaries using patterns
 - Concise explanations
 - Meaningful (semantically)
- Integrating sampling into provenance capture process
 - $_{\circ}$ Unbiased
 - Computing representative patterns
 - Calculating close enough approximate completeness of patterns



• What are patterns?

 $r_1(N, shared, I, apt, E, P)$ (T,F)



• What are patterns?

 $r_1(N, shared, I, apt, E, P)$ (T,F)

r_1 (central place, shared, 8403, apt, east, 130)	(T , F)
r_1 (near SpaceNeedle, shared, 8575, apt, lower, 40)	(T , F)
r_1 (central place, shared, 8403, apt, east, 40)	(T , F)
r_1 (near SpaceNeedle, shared, 8575, apt, lower, 350)	(T,F)
r_1 (central place, shared, 8403, apt, east, 350)	(T , F)
r_1 (central place, shared, 8403, apt, east, 45)	(T , F)

....



• What are patterns?

 $r_1(N, shared, I, apt, E, P)$ (T,F)

r_1 (central place, shared, 8403, apt, east, 130)	(T,F)
r_1 (near SpaceNeedle, shared, 8575, apt, lower, 40)	(T , F)
r_1 (central place, shared, 8403, apt, east, 40)	(T , F)
r_1 (near SpaceNeedle, shared, 8575, apt, lower, 350)	(T , F)
r_1 (central place, shared, 8403, apt, east, 350)	(T,F)
r_1 (central place, shared, 8403, apt, east, 45)	(T , F)

....



• What are provenance summaries?





- Quality metrics
 - ^o Completeness (cp): fraction of provenance covered by a pattern





- Quality metrics
 - ^o Completeness (cp): fraction of provenance covered by a pattern
 - Informativeness (info): degree of new information from a pattern





- How to compute summaries
 - Heuristic using a sample of why-not provenance





• How to compute summaries



Summarization process





- Sampling why-not provenance
 - Generating a sample that is equivalent to uniform random sample
 - Without computing full why-not provenance
 - Batch sampling: generating a query that returns an unbiased sample





- Generating patterns
 - LCA (Lowest Common Ancestor)



Interpretable and Informative Explanations of Outcomes [Gebaly et al. PVLDB 2014]



• Measuring quality and selecting top-k patterns



C: completeness I: Informativeness





- Performance of computing summaries
- Quality of summaries
- Comparison with other approaches
- Datasets
 - 4 real-world datasets
 - TPCH
- Queries
 - Single rule through multiple rules
 - Negation and comparisons

Approximate summaries for why and why-not provenance (extended version) [Lee et al. https://arxiv.org/abs/2002.00084 2020]



• Computing summaries for ~10⁵⁰ derivations





Generating high-quality summaries



[Completeness comparison]

[Quality metric error caused by sampling]